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**Final  
CSE 131B  
Spring 2004**

**Page 1** \_\_\_\_\_ (25 points)

**Page 2** \_\_\_\_\_ (24 points)

**Page 3** \_\_\_\_\_ (32 points)

**Page 4** \_\_\_\_\_ (24 points)

**Page 5** \_\_\_\_\_ (28 points)

**Page 6** \_\_\_\_\_ (26 points)

**Page 7** \_\_\_\_\_ (22 points)

**Page 8** \_\_\_\_\_ (23 points)

**Subtotal** \_\_\_\_\_ (204 points)

**Page 9** \_\_\_\_\_ (10 points)

**Extra Credit**

**Total** \_\_\_\_\_

**1a.** Consider the following C/C++ program:

```
#include <stdio.h>

int main()
{
    int y = 1;
    int x[5];
    int i;

    for ( i = 0; i <= 5; ++i )
        x[i] = i;

    printf( "%d\n", y );      /* Print the value of y */
    printf( "%p\n", &y );    /* Print the address of y */

    return 0;
}
```

Explain why the first printf() statement prints the value 5 when compiled with some compiler on some architecture (it does on SPARC with cc and CC)? (5 points)

Explain why the first printf() statement prints the value 1 when compiled with the same compiler on the same architecture as above but with the second printf() statement removed? (5 points)

**1b.** C/C++ structs are assignable only if they are equivalent types. Two structs that are not equivalent cannot be assigned even with a cast. Why? (5 points)

How can you get around this limitation? Give a specific example. (5 points)

Why is this OKAY? (5 points)

[Hint: Discuss types of equivalence used in your answers]

```
int
leapyear( unsigned int year )
{
    if ( (year % 400 == 0) || ((year % 4 == 0) && (year % 100 != 0)) )
        return 1;
    else
        return 0;
}
```

```
int a[3][2];
```

```
VAR a : ARRAY 3, 2 OF INTEGER
```

(4 points)

a[2,0]

low memory high memory

3. Why is passing and returning references or pointers to structs/objects as arguments and return values usually recommended over passing/returning structs/objects by value in most languages? (4 points)

What gets printed? (28 points)

```
VAR x : INTEGER;
VAR y : BOOLEAN;

PROCEDURE fool( a : BOOLEAN; VAR b : INTEGER ) : INTEGER;
  VAR i : INTEGER;
  VAR j : BOOLEAN;
BEGIN
  i := b;
  b := 77;
  j := a;
  a := FALSE;

  OUTPUT x, " ", y;      _____
  OUTPUT a, " ", b;      _____
  OUTPUT i, " ", j;      _____

  RETURN i;
END fool;

PROCEDURE foo( a : INTEGER; VAR b : BOOLEAN );
  VAR i, j : INTEGER;
BEGIN
  i := a;
  a := 66;
  j := fool( b, a );
  b := TRUE;

  OUTPUT x, " ", y;      _____
  OUTPUT a, " ", b;      _____
  OUTPUT i, " ", j;      _____
END foo;

BEGIN
  x := 55;
  y := FALSE;

  foo( x, y );

  OUTPUT x, " ", y;      _____
END.
```

4. Identify where each of the following program parts live in the Java runtime environment as discussed in class. (12 points)

public class Foo {		
private static Foo a;	a	_____
private int b;	b	_____
public Foo() {	Foo()	_____
a = this;	this	_____
++b;		
}		
	main()	_____
public static void main( String[] args ) {	args	_____
Foo c = new Foo();	c	_____
int d;	d	_____
c = new Foo();	where c is pointing	_____
c.method( d );		
}	method()	_____
private void method( int e ) {	e	_____
int f;	f	_____
f = e;		
}		
}		

Using the Right-Left rule (which follows the operator precedence rules) write the definition of a variable named foo that is a pointer to an array of 9 elements where each element is a pointer to a function that takes a pointer to a struct Pub as the single parameter and returns a pointer to a 7x19 2-D array where each element is a pointer to a struct Fubar. (6 points)

Regarding type checking, reference (VAR) parameters require the actual arguments to be \_\_\_\_\_  
 and \_\_\_\_\_ to the formal parameter type while value parameters require the actual arguments  
 to be \_\_\_\_\_ to the formal parameter type. (6 points)

5. Given the following program, order the printf() lines so that the values that are printed when run on a Sun SPARC Unix system are displayed from smallest value to largest value and label the corresponding C/C++ Runtime area. (20 points)

void foo( int, int * ); /* Function Prototype */		
int a;		
int main( int argc, char *argv[] ) {	smallest value (low memory)	Runtime Area
static int b = 420;		
int c = 404;		
foo( argc, &c );		
/* 1 */ (void) printf( "a --> %p\n", &a );		
/* 2 */ (void) printf( "b --> %p\n", &b );		
/* 3 */ (void) printf( "c --> %p\n", &c );		
/* 4 */ (void) printf( "argc --> %p\n", &argc );		
/* 5 */ (void) printf( "malloc --> %p\n", malloc(50) );		
/* 6 */ (void) printf( "foo --> %p\n", foo );		
}		
void foo( int d, int *e ) {		
int f = 911;		
int g;		
/* 7 */ (void) printf( "d --> %p\n", &d );		
/* 8 */ (void) printf( "e --> %p\n", &e );		
/* 9 */ (void) printf( "f --> %p\n", &f );		
/* 10 */ (void) printf( "g --> %p\n", &g );	largest value (high memory)	
}		

Why do compilers typically allocate space for arguments in the Runtime Stack even when they pass them in registers? (4 points)

Object-oriented languages allow operator, function, and constructor overloading. In these languages, the function name is not always a unique identifier, since you can have multiple related definitions. For lookup purposes, the compiler must construct a distinct identifier for each function. Sometimes, such overloaded functions will have different return types as well. How would you create distinct identifiers for such functions? (4 points)

6. Given the following Oberon program, emit the **unoptimized** SPARC assembly language code that should be generated for the two procedures `foo1()` and `foo()`. Assume no optimizations – treat each instruction separately without any knowledge of any previously computed/loaded/stored values that may still be in a register from a previous instruction. Draw a line between each group of assembly language instructions that represent the emitted code generated for each instruction and label them with the instruction number. All local variables must be allocated and accessed on the Stack (do not map them directly into a local register). (26 points)

```

PROCEDURE foo1( VAR x : INTEGER; VAR y : INTEGER ) : INTEGER;
  VAR i, j : INTEGER;          (**** Local Stack variables - Do not initialize to zero ****)
BEGIN
  j := y + 5;                  (* 1 *)
  i := x;                      (* 2 *)
  y := j - 10;                  (* 3 *)
  RETURN x;                     (* 4 *)
END foo1;

PROCEDURE foo( a : INTEGER; VAR b : INTEGER );
  VAR i, j : INTEGER;          (**** Local Stack variables - Do not initialize to zero ****)
BEGIN
  i := foo1( a, b );           (* 5 *)
  b := foo1( i, j );            (* 6 *)  (* Don't worry about any local variables not initialized *)
END foo;

BEGIN
  (* ... *)
END.

.global foo, foo1, main

.section ".text"

foo1:                                foo:

```

7. Show the memory layout of the following C struct/record definition taking into consideration the SPARC data type memory alignment restrictions discussed in class. Fill bytes in memory with the appropriate struct/record member/field name. For example, if member/field name `p` takes 4 bytes, you will have 4 `p`'s in the appropriate memory locations. If the member/field is an array, use the name followed by the index number. For example, some number of `p0`'s, `p1`'s, `p2`'s, etc. Place an X in any bytes of padding. Structs are padded so the total size is evenly divisible by the most strict alignment requirement of its members. (11 points)

```

struct foo {
    char    a;
    short   b;
    double  c;
    short   d[3];
    int     e;
}

struct foo fubar;

```

fubar:

low memory


Generate SPARC assembly code for the 3 labeled statements below. Assume the two local variables have been allocated on the Stack with an appropriate `save` instruction. Draw a line between and label the group of translated instructions for each labeled line. (11 points)

```

PROCEDURE P();
    VAR i : INTEGER;
    VAR ptr : POINTER TO INTEGER;
BEGIN
    NEW( ptr );                (* 1 *)

    ptr^ := 420;               (* 2 *)

    i := ptr^;                 (* 3 *)
END P;

BEGIN
END.

```



8. Give the order of the phases of compilation in a typical compiler as discussed in class (8 points)

A – Machine-specific code improvement (optional)	B – Scanner (lexical analysis)
C – Parser (Semantic analysis/intermediate code gen.)	D – Parser (syntax analysis)
E – Machine-independent code improvement (optional)	F – Target code generation
G – Source language (for example, C)	H – Target language (for ex., assembly)

\_\_\_\_\_ -> \_\_\_\_\_ -> \_\_\_\_\_ -> \_\_\_\_\_ -> \_\_\_\_\_ -> \_\_\_\_\_ -> \_\_\_\_\_ -> \_\_\_\_\_

Give the order of the typical C compilation stages and on to actual execution as discussed in class (8 points)

B – loader	G – cpp (C preprocessor)
D – ld (Linkage Editor)	E – as (assembler)
C – exe/a.out (executable image)	H – Source file
F – ccomp (C compiler)	A – Program Execution

gcc \_\_\_\_\_ -> \_\_\_\_\_ -> \_\_\_\_\_ -> \_\_\_\_\_ -> \_\_\_\_\_ -> \_\_\_\_\_ -> \_\_\_\_\_ -> \_\_\_\_\_

How did you implement the OUTPUT function in your Project 2? (5000 points)

Tell me something you learned in this class that is extremely valuable and that you think you will be able to use for the rest of your programming/computer science career. (2 points)

## 9. Extra Credit (10 points)

Given the following ANSI/ISO C/C++ variable definitions, identify which expressions will produce a static semantic compiler error.

A) No compiler error

B) Compiler error

```
int
main()
{
    char * s = "CSE 131B Rocks!";

    char * p1 = &s[4];
    const char * p2 = s;
    char * const p3 = s + 7;
    const char * const p4 = &(s + 9);

    p1 = s; _____
    *p1 = 'A'; _____
    p2 = s; _____
    *p2 = 'A'; _____
    p3 = s; _____
    *p3 = 'A'; _____
    p4 = s; _____
    *p4 = 'A'; _____
    *&p1[9] = *(s + 1); _____
    ((char *) p4) = (char *) p3; _____
}
```

Note: cc, CC, and g++ report these as errors; gcc reports these as warnings! :-/

## Scratch Paper

## Scratch Paper